

- Applications of Monte Carlo simulations.
- Random number generators.

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- Coin toss: the only prediction about the outcome: 50% of the events will end up as tail being up.
- Dice: In a large number of throws, the probability of getting a given face is $\frac{1}{6}$.

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- Often need a program that generates a random variable with a given probability distribution.

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- This is because computers have only a limited number of bits to represent a number.
- It implies that no matter which pseudo random number generator you use – it will always repeat itself (period of the generator).

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- Fast.

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- Feedback shift register methods

Generates a pseudo random sequence of numbers $\{x_1, x_2, \dots, x_k\}$ of length M over the interval $[0, M - 1]$:

$$x_i = \text{mod } (ax_{i-1} + c, M) = \text{remainder} \left(\frac{ax_{i-1} + c}{M} \right)$$

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$a = 4, c = 1, M = 9, x_1 = 3$ leads to the following sequence:

$$x_2 = 4$$

$$x_3 = 8$$

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interval: $0 - 8$ i.e. $[0, M - 1]$

Period: 9 i.e. M numbers (then repeat).

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- For $c = 0$ called "Multiplicative congruential generator".

- Scale results from x_i on $[0, M - 1]$ to y_i on $[0, 1]$.

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- Scale results from y_i on $[0, 1]$ to z_i on $[A, B]$

$$z_i = A + (B - A) * y_i$$

Simple shift register where the vacated bit is filled with the exclusive-or of two other bits in the shift register.

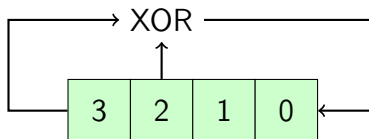
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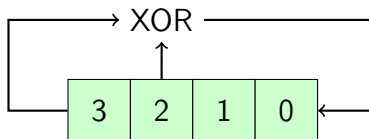
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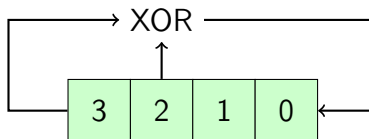


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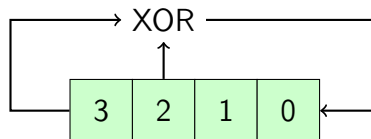


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4 bit shift-register pseudorandom number generator:



- Bits 3 and 2 are combined by exclusive-or.
- The register is shifted 1 step to the left.
- The result of the exclusive-or is entered into bit 0.

4bit shift register PRNG

Here is the pattern of bits, starting with 0001:

0001

0010

0100

1001

0011

0110

1101

1010

0101

1011

0111

1111

1110

1100

1000

0001

- Most commonly used is Mersenne Twister (which is a generalized feedback shift register method).

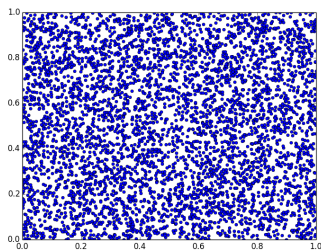
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- Implemented in numpy.

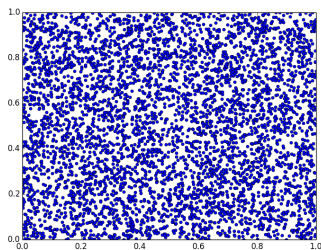
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- 2D plot, x_i and y_i from two random sequences (parking lot test).



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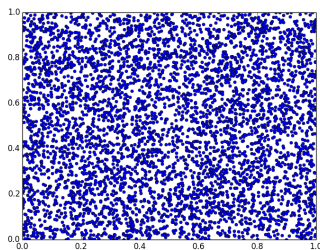
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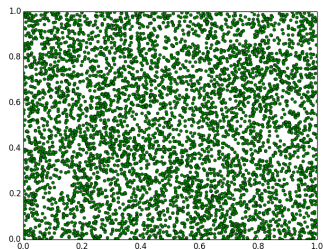
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- Plot 3D figure (x_i, y_i, z_i)
- Plot correlation (x_i, x_{i+k})



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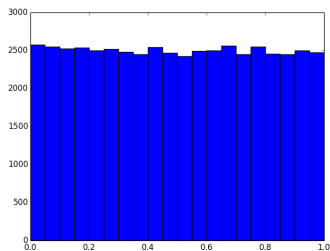
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- Near neighbour correlation:

$$\frac{1}{N} \sum_{i=1}^N x_i x_{i+k} \approx \frac{1}{4}$$

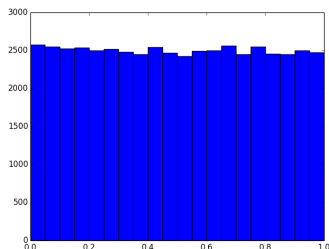
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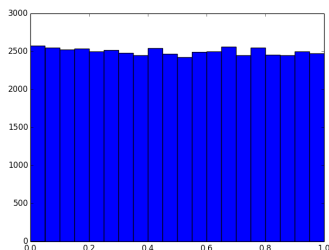


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- 4th moment: (50000 random numbers)

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- near neighbor correlation: (50000 random numbers)
= 0.2478

Good test suites exist – TestU01 – which can be used to uncover problems in random number generators.

Dont try to invent your own random number generator – unless you know what you are doing. This is very tricky business!!!